

An International Open Access Journal
UGC and ISSN Approved | ISSN: 2349-5162

INTERNATIONAL JOURNAL
OF EMERGING TECHNOLOGIES
AND INNOVATIVE RESEARCH

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INTERNATIONAL JOURNAL OF EMERGING
TECHNOLOGIES AND INNOVATIVE RESEARCH

International Peer Reviewed, Open Access Journal
ISSN: 2349-5162 | Impact Factor: 5.87

UGC and ISSN Approved Journals.

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A note on Linear codes associated to Schubert varieties

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Abstract:

We consider the linear code associated with Schubert sub variety of Grassmannian's. In this review article we have studied the basic notions of Schubert code $C_\alpha(l, m)$. We have discussed the some known results and examples of Schubert code.

Key Words : Linear code, Grassmannian, MDC, Schubert

1.Introduction:

Let q be the power of fixed prime and l and m be the positive integers with $l \leq m$. Let F_q denote the finite field with q elements and F_q^m be the vector space of dimension m over F_q . Let $\begin{bmatrix} m \\ l \end{bmatrix}_q$ denote Gaussian binomial coefficient and $G_{l,m}$ denote the Grassmannian of all l -planes of F_q^m . It is also known that the Grassmannian $G_{l,m}$ embeds into projective space $P(F_q)^{\binom{m}{l}-1} = P(F_q)^{\binom{m}{m-l}-1}$ via plucker embedding. The image under this mapping is projective algebraic variety. To every projective space one can associate a linear code (example in [1]). The linear code corresponding to $G_{l,m}$ is known as Grassmannian code and it is denoted by $C(l, m)$. The Grassmannian code were introduced by C. T. Ryan in [2,3] for binary case. D Yu Nogin in [4] studied the linear code $C(l, m)$ associated to the Grassmannian $G_{l,m}$ over finite field and verified that $C(l, m)$ is an $[n, k, d]_q$ code where

$$n = \begin{bmatrix} m \\ l \end{bmatrix}_q = \frac{(q^m-1)(q^m-q)\dots(q^m-q^{l-1})}{(q^l-1)(q^l-q)\dots(q^l-q^{l-1})}, k = \binom{m}{l} \text{ and } d = q^{l(m-l)}$$

For fix integers k, n with $1 \leq k \leq n$ and prime q . Let C be linear $[n, k]_q$ code i.e C is k -dimensional subspace of F_q^n . Given for any $x = (x_1, x_2, \dots, x_n) \in F_q^n$, define $\text{sup}(x) = \{i: x_i \neq 0\}$ and $\|x\| =$

$|\text{sup}(x)|$ denote support and Hamming norm of x , likewise for $D \subseteq F_q^n$ $\text{sup}(D) = \{i: x_i \neq 0 \text{ for some } x = (x_1, x_2, \dots, x_n) \in D\}$ and $\|D\| = |\text{sup}(D)|$ denote support and hamming norm of D . The minimum distance or hamming weight of C is defined by $d(C) := \min\{\|x\|: x \in C \text{ with } x \neq 0\}$.

For some positive integer r , the r^{th} higher weight or r^{th} generalized hamming weight is denoted by $d_r = d_r(C)$ of the code C defined by

$$d_r(C) = \min\{\|D\|: D \text{ is subspace of } C \text{ and } \dim(D) = r\}$$

The set $\{d_r(C) : 1 \leq r \leq k\}$ is complete weight hierarchy of the code C . For linear code C it is very interesting and difficult to determine complete weight hierarchy.

Let $\alpha = (\alpha_1, \alpha_2, \dots, \alpha_l)$ be strictly increasing sequence of positive integers satisfying

$$\alpha_1 \leq \alpha_2 \leq \dots \leq \alpha_l \leq m \text{ and } \Omega_\alpha(l, m) \text{ be corresponding Schubert variety in Grassmannian } G_{l,m}$$

Schubert varieties are sub-varieties of Grassmannian. Likewise Grassmannian varieties, the Schubert variety may be considered as subset of $P^{\binom{m}{l}-1}$. The linear code corresponding to Schubert variety $\Omega_\alpha(l, m)$ is called Schubert code and this code is denoted by $C_\alpha(l, m)$. Ghorpade-Tsfasman [5] had proved the minimum distance conjecture for Schubert code corresponding to case when $\delta(\alpha) = l(m-l) - 1$ and found length and dimension of Schubert code in general. In [5] Ghorpade-Tsfasman proved that Schubert code $C_\alpha(l, m)$ is $[n_\alpha, k_\alpha]$ code where $n_\alpha = \sum_{\beta \leq \alpha} q^{\delta(\beta)}$, $k_\alpha = \det_{1 \leq i, j \leq l} \left(\binom{\alpha_j - j + 1}{i - j + 1} \right)$ where sum is run over all l -tuples $\beta = (\beta_1, \beta_2, \dots, \beta_l)$ of integers satisfying $1 \leq \beta_1 \leq \beta_2 \leq \dots \leq \beta_l \leq m$ and $\beta_i \leq \alpha_i$ for $i = 1, 2, 3, \dots, l$ and $\delta(\beta) = \sum_{i=1}^l (\beta_i - i)$. Alternating proof of MDC for Schubert code was given in [6] and [7].

2. Preliminaries :

The Grassmannian of all l -planes of F_q^m is given by,

$$G_{l,m} := \{L \subseteq F_q^m : L \text{ is subspace of } F_q^m \text{ and } \dim L = l\}.$$

Due to plucker embedding of $G_{l,m}$ into projective space $P^{\binom{m}{l}-1}$, Choose matrix A_L whose rows forms basis for L . The order of A_L will be $l \times m$

matrix with rank l . In A_L we have $\binom{m}{l}$ minors of order l . By fixing some ordering of these minors and map L onto $\binom{m}{l}$ tuples of minors of A_L of size l which is required embedding of $G_{l,m}$ into projective space $P^{\binom{m}{l}-1}$.

Let $X = (X_{ij})$ be $l \times m$ indeterminate matrix over F_q . Let l -multiset $I \subset \{1, 2, \dots, m\}$ denote $l \times l$ minor of X corresponding to the columns indexed by I by $\det_I(X)$ and $F_q[X]_l$ be a vector space over F_q spanning by minors of $\det_I(X)$. Then for any L in $G_{l,m}$, the $l \times m$ matrix A_L is matrix whose rows span L .

Consider the evaluation map $Ev : F_q[X]_l \rightarrow F_q^{\binom{m}{l}}$ defined by $f = \sum_I \alpha_I \det_I(X) \rightarrow (f(A_L))_{L \in G_{l,m}}$ and $f(A_L) = \sum_I \alpha_I \det_I(A_L)$ i.e. $f(A_L)$ is evaluation of f at A_L . The image of this is called Grassmann code $\mathcal{C}(l, m)$. The image $Ev(f)$ of f denoted by c_f is codeword corresponding to f .

Let $\alpha = (\alpha_1, \alpha_2, \dots, \alpha_l)$ be the sequence of strictly increasing positive integers and $A_1 \subset A_2 \dots \subset A_l$ be the sequence of subspaces of F_q^m with $\dim A_i = \alpha_i \forall i$.

Let $\Omega_\alpha(l, m) = \{W \in G_{l,m} : \dim(W \cap A_i) \geq i, \forall i\}$ be the Schubert variety in $G_{l,m}$ corresponding to sequence α . If the above evaluation map is restricted to $\Omega_\alpha(l, m)$ then it will be Schubert code $C_\alpha(l, m)$.

The Schubert variety $\Omega_\alpha(l, m)$ only depends on the sequence α not on corresponding sequence $A_1 \subset A_2 \dots \subset A_l$ of subspaces. If $B_1 \subset B_2 \dots \subset B_l$ is another sequence of subspaces of F_q^m with $\dim B_i = \alpha_i$ for every i . If $\Omega_\alpha(A_1, A_2, \dots, A_l, m)$ denotes the Schubert varieties corresponding to the sequence of subspaces $A_1 \subset A_2 \dots \subset A_l$ and $\Omega_\alpha(B_1, B_2, \dots, B_l, m)$ denotes Schubert variety corresponding to sequence of subspaces $B_1 \subset B_2 \dots \subset B_l$ then due to [1] these corresponding Schubert codes are equivalent.

3. Basic Notions and some known facts about Schubert Code :

For fix integer l, m with $1 \leq l \leq m$. Let $I(l, m)$ be the indexing set with partial order \leq for any $\beta = (\beta_1, \beta_2, \dots, \beta_l) \in I(l, m)$.

$$\text{Let } \delta_\beta := \sum_{i=1}^l (\beta_i - i) = \beta_1 + \beta_2 + \dots + \beta_l - \frac{l(l+1)}{2}$$

For fix some $\alpha \in I(l, m)$ then $C_\alpha(l, m)$ be the corresponding linear code and length of Schubert code [9] is n_α given by

Thm 1 (Length of Schubert Code)[9]The number of F_q rational points of Ω_α which also length n_α of $C_\alpha(l, m)$ is given by $n_\alpha = \sum_{\beta \leq \alpha} q^{\delta_\beta}$ where the sum is over all $\beta \in I(l, m)$ satisfying $\beta \leq \alpha$

With the previous notations the explicit formula for the dimension of Schubert code $C_\alpha(l, m)$ is K_α is given by [9] in following result .

Thm 2(Dimension of Schubert Code)

The dimension of Schubert code $C_\alpha(l, m)$ is given by $l \times l$ determinant and given due to [9] by

$$K_\alpha = \det_{1 \leq i, j \leq l} \left(\binom{\alpha_j - j + 1}{i - j + 1} \right) = \begin{vmatrix} \binom{\alpha_1}{1} & 1 & 0 & \dots & 0 \\ \binom{\alpha_1}{2} & \binom{\alpha_2 - 1}{1} & 1 & \dots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ \binom{\alpha_1}{l} & \binom{\alpha_2 - 1}{l-1} & \binom{\alpha_3 - 2}{l-2} & \dots & \binom{\alpha_l - l + 1}{1} \end{vmatrix}$$

Thm 3[9]: The dimension k_α of q -ary Schubert code $C_\alpha(l, m)$ is independent of q and is related to the length $n_\alpha = n_\alpha(q)$ of $C_\alpha(l, m)$ by the formula $\lim_{q \rightarrow \infty} \frac{n_\alpha(q)}{q} = k_\alpha$

4.Examples of Schubert Code :

Example 4.1

Let $\alpha = (2, 4)$ then $C_\alpha(2, 4)$ is corresponding Schubert code then dimension of $C_\alpha(2, 4)$ is 5

Example 4.2

Let $\alpha = (3, 5)$ then $C_\alpha(3, 5)$ is corresponding Schubert code and dimension of $C_\alpha(3, 5)$ is 9

Acknowledgement : Author is very much thankful to UGC for providing financial support in the form of Minor research project under XII plan .

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